

GATE PSUs

State Engg. Exams

MADE EASY
WORKBOOK 2024



**Detailed Explanations of
Try Yourself *Questions***

Computer Science & IT
Programming
and Data Structures



1

Programming



Detailed Explanation of Try Yourself Questions

T1 : Solution

[$O(n^2)$]

$A(n)$

```
{   for (i = 1 to n)
    {   if (n mod i == 0)
        {   for (j = 1 to n)
            printf(j)
        }
    }
}
```

$\left. \begin{array}{l} \left. \begin{array}{l} \text{if } (n \bmod i == 0) \\ \text{for } (j = 1 \text{ to } n) \\ \text{printf}(j) \end{array} \right\} = O(n/2) = O(n) \end{array} \right\} = O(n)$

Time complexity = $O(n) \times O(n) = O(n^2)$

T2 : Solution

[$O(1)$]

main()

```
{   int i = 3;
    switch (i)
    {   default : printf("zero")
        Case 1 : printf("one")
            break
        Case 2 : printf("two")
            break
        Case 3 : printf("three")
            break
    }
}
```

Since $i = 3$ so switch (3) will go to case 3 and run the program only one time.
So time complexity = $O(1)$.

T3 : Solution

1. `Const int *P;`
declare P as pointer to const integer.
2. `int * const P;`
declare P as constant pointer to integer

T4 : Solution

- (i) `Char (*( *x ( ) ) [ ] ) ();`
declare x as a function returning pointer to array of pointer to function returning char.
- (ii) `Char (*( *x [3] ) ( ) [5];`
declare x as array 3 of pointer to function returning pointer to array 5 of char.
- (iii) `Void (*b*int, void (*f)(int)) (int);`
Syntac error
- (iv) `Void (*ptr)(int (*)[2], int(*) (void));`
Syntax error

T5 : Solution

(b)

`Char \0`

if (0) `∴ ∴ ∴ Printf(“% S”, a) = Null = 0`

So condition false

So answer is else part string is not empty.

T6 : Solution

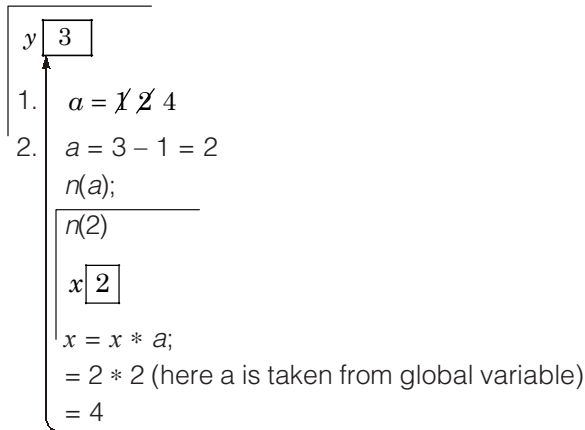
(a)

Since variable d of integer type is static so memory is allocated to it compile time only and same memory is used every time. Therefore, every time old value of d (which is update in previous iteration) is used.

So, output is 312213444.

T7 : Solution

(d)

 a globally initialize.1. $m(3)$  $\text{Printf}(4); = 4$

3. $\text{Printf}(a) = 4$ since dynamic scoping is used. So, take value of inner variable 'a'.
So answer will be 4, 4.

T8 : Solution

(c)

Take random value of X and Y i.e., $X = 5$ and $Y = 3$.Initially $X = 5, Y = 3, \text{res} = 1, a = X$ and $b = Y$ **Option (a):** $X^Y = a^b$

$$X^Y = a^b \equiv 5^3 = 5^3 \equiv 125 = 125$$

After iteration 1

$$\text{res} = 5; a = 5; b = 2; X = 5; Y = 3$$

$$X^Y = a^b \equiv 5^3 \neq 5^2 \equiv 125 \neq 25$$

So, case fail. Option (a) cannot be answer.

Option (b): $(\text{res} * a)^Y = (\text{res} * X)^b$

$$(1 \times 5)^3 = (1 \times 5)^3 \equiv 125 = 125$$

After iteration 1

$$\text{res} = 5; a = 5; b = 2; X = 5; Y = 3$$

$$(\text{res} * a)^Y = (\text{res} * X)^b \equiv (5 \times 5)^3 = (5 \times 5)^2$$

15625 $\neq$ 625 So, case fail. Option (b) cannot be answer.**Option (d):** $X^Y = (\text{res} * a)^b$

$$5^3 = (1 \times 5)^3 \equiv 125 = 125$$

After iteration 1

$$\text{res} = 5; a = 5; b = 2; X = 5; Y = 3$$

$$X^Y = (\text{res} * a)^b \equiv 5^3 = (5 \times 5)^2$$

125 $\neq$ 625 So, case fail.

Option (d) cannot be answer.

Option (c):

$$X^Y = \text{res} * a^b$$

$$5^3 = 1 \times 5^3 \equiv 125 = 125$$

After iteration 1

$$\text{res} = 5; a = 5; b = 2; X = 5; Y = 3$$

$$X^Y = \text{res} * a^b \equiv 5^3 = 5 \times 5^2 \equiv 125 = 125$$

After iteration 2

$$\text{res} = 25; a = 5; b = 1; X = 5; Y = 3$$

$$X^Y = \text{res} * a^b \equiv 5^3 = 25 \times 5^1 \equiv 125 = 125$$

So, all cases are passes.

So option (c) will be the answer.

T9 : Solution

a	3	5	2	6	4
---	---	---	---	---	---

1. $f(a, 5)$ is a function contain 2 parameter one contain starting address of array and second parameter tell number of element in the array.
2. Every time 'n' value compare with 1 when it is less than equal to 1 return 0 and stop the program otherwise continue with recursive function call.

1. $f(a, 5)$

*P = a; P pointed to same address pointed by 'a'.

n = 5; n value greater than 1.

So, $\max(f(P + 1, 5 - 1), 3 - 5)$; or

$\max(f(P + 1, 4) - 2)$;

2. $f(P + 1, 4)$

*P = P + 1; P is pointed to next element of array i.e., 5.

n = 4; n value greater than 1.

So, $\max(\max(f(P + 1, 4 - 1), 5 - 2), -2)$ or

$\max(\max(f(P + 1) 3), 3), -2)$

3. $f(P + 1, 3)$

*P = P + 1; P is pointed to next element of array i.e., 2.

n = 3; n value greater than 1.

So, $\max(\max(\max(f(P + 1, 3 - 1), 2 - 6), 3) - 2)$ or $\max(\max(\max(f(P + 1, 2), -4), 3) - 2)$;

4. $f(P + 1, 2)$;

*P = P + 1; P is pointed to next element of array i.e., 6.

n = 2; n value greater than 1.

So, $\max(\max(\max(\max(f(P + 1, 2 - 1), 2), -4), 3), -2)$ or

$\max(\max(\max(\max(f(P + 1, 1), 2), -4), 3), -2)$

5. $f(P + 1, 1)$;

*P = P + 1; P is pointed to next element of array i.e., 4.

n = 1; n value equal to 1 so, return 0.

So $\max(\max(\max(\max(0, 2), -4), 3), -2)$

$\max(\max(\max(2, -4), 3), -2)$

$\max(\max(2, 3), -2)$

$\max(3, -2) = 3$

So the value printed by given code is 3.



2

Linked List, Stack, Queue and Hashing

T1 : Solution

Implementation of stack using single linklist:

Inserting sequence: 1, 2, 3, 4, 5, 6

Insertion take $O(1)$ time

Link list representation:

1. $\rightarrow$

1	/
---	---

2. $\rightarrow$

2	
---	--

 $\rightarrow$

1	/
---	---

3. $\rightarrow$

3	
---	--

 $\rightarrow$

2	
---	--

 $\rightarrow$

1	/
---	---

4.

5.

6. Insertion takes $O(1)$ time.

Deletion in stack (Pop)

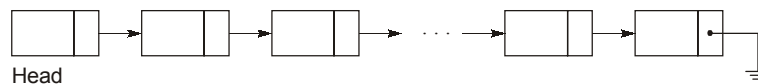
Remove top element every time so $O(1)$

Deletion in linklist

Remove 1st node every time with making second node to head.

6
5
4
3
2
1

T2 : Solution



enqueue operation takes $O(1)$ time

dequeue operation takes $O(n)$ time [visits last node]

T3 : Solution

(d)

PUSH (S, P, Q, T_i , x)

{

if $\left(T_i == \left(\frac{P}{Q} \times (i+1) - 1\right)\right)$

{

printf ("stack overflow");

exit (1);

}

else

T_i++ ;

$S[T_i] = x$;

}

$T_i == \left(\frac{P}{Q} \times (i+1) - 1\right)$ indicate the last location of the array is already filled. So overflow occur.

T4 : Solution

(a)

Number of push operations = $n(\text{insert}) + m(\text{delete}) = n + m$

So, $n + m \leq x$ but there are maximum $2n$ insert operations so $n + m \leq x \leq 2n$... (1)

Number of pop operations = $n + m$

But there are $2m$ delete operations which are less than no. of pop operations, hence

$$2m \leq n + m \quad \dots (2)$$

From (1) and (2): $n + m \leq x \leq 2n$ and $2m \leq n + m$

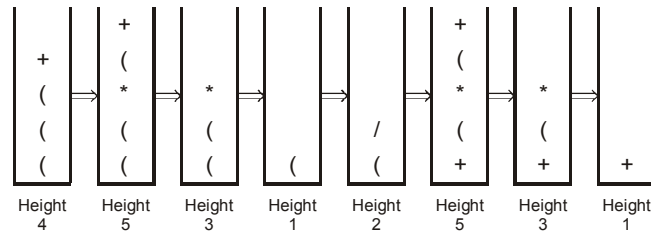
T5 : Solution

(22079)

$$\begin{aligned} \text{Formula to find location of } a[20][20][30] &= 10 + \{[(20-1)(30-1)(40-1)] + (20-1)(30-1) + (30-1)\} \\ &= [10 + (19 \times 29 \times 39) + (19 \times 29) + (29)] \\ &= 10 + 21489 + 551 + 29 \\ &= 10 + 22069 \\ &= 22079 \end{aligned}$$

T6 : Solution

(5)



(Unique heights are 1, 2, 3, 4, 5) where 1, 3, 5 are repeated two times each. Maximum size of stack is 5.

T7 : Solution

(0.7324)

Expected number of probes in a unsuccessful = $\frac{1}{(1-\alpha)}$

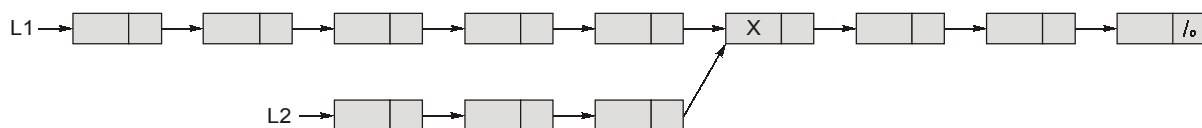
$$\begin{aligned}\frac{1}{1-\alpha} &= 3 \\ 1 &= 3(1-\alpha) \\ 1 &= 3-3\alpha \\ -2 &= -3\alpha \\ \alpha &= \frac{2}{3}\end{aligned}$$

Expected number of probes in a unsuccessful = $1/\alpha \log_e 1/(1-\alpha)$

$$\frac{3}{2} \log_e 3 = 0.7324$$

T8 : Solution

(b)



We need to traverse both the linked list of size m and n .

So it will take $O(m + n)$.

T9 : Solution

(b)

By using BSF (Breadth First Search) traversal we can set the twin pointer in each entry in each adjacency list.

So it will take $\Theta(m + n)$ times (since adjacency list are using).

