

# CLASS TEST

S.No. : 04 LS1\_EE\_T\_230819

Digital Electronics



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# CLASS TEST 2019-2020

## ELECTRICAL ENGINEERING

### Digital Electronics

Date of Test : 23/08/2019

#### Answer Key

- |        |         |         |         |         |
|--------|---------|---------|---------|---------|
| 1. (d) | 7. (b)  | 13. (d) | 19. (c) | 25. (a) |
| 2. (b) | 8. (*)  | 14. (c) | 20. (b) | 26. (a) |
| 3. (c) | 9. (d)  | 15. (c) | 21. (d) | 27. (a) |
| 4. (a) | 10. (a) | 16. (a) | 22. (a) | 28. (c) |
| 5. (d) | 11. (a) | 17. (a) | 23. (a) | 29. (a) |
| 6. (b) | 12. (a) | 18. (a) | 24. (b) | 30. (d) |

## DETAILED EXPLANATIONS

1. (d)

Given 12 bit content in excess 3 form

$$1000\ 1001\ 0111 = 897 - 333 = 564$$

3. (c)

Dual of Ex-OR is Ex-NOR. Also, complement of Ex-OR is Ex-NOR which is dual of Ex-OR.

4. (a)

CLK	State	Serial in
X	0 0 1 1	0
1	0 0 0 1	1
2	1 0 0 0	

5. (d)

Output of stage 1  $X \oplus X = 0$ Output of 2nd stage  $1 \oplus 0 = 1$ Output  $Y = 1 \oplus 0 = 1$ 

6. (b)

Minimum number of FF required,

$$2^n \geq \text{Mod no.}$$

$$\therefore 2^n \geq 128$$

$$n = 7$$

8. (\*)

$$\overline{(A \cdot \bar{B} + \bar{C})D + \bar{E}}$$

$$\overline{(A \cdot \bar{B} + \bar{C})D \cdot E}$$

$$\overline{((\overline{A \cdot \bar{B} + \bar{C}}) + \bar{D})E}$$

$$\overline{((\overline{A \cdot \bar{B} \cdot C}) + \bar{D})E}$$

$$((\bar{A} + B) \cdot C + \bar{D})E$$

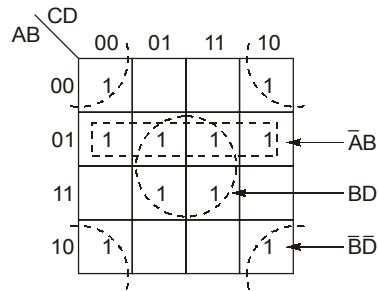
9. (d)

$$(1010) = 10_{10} \text{ gives } 3 \text{ V}$$

$$\therefore \text{Step size} = \frac{3\text{V}}{10} = 0.3 \text{ V}$$

$$\text{Output voltage for } 0100 = (4_{10}) \text{ is } 4 \times 0.3 = 1.2 \text{ V}$$

10. (a)  
k map



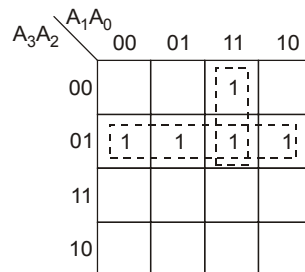
∴

$$f = BD + \bar{A}B + \bar{B}\bar{D}$$

11. (a)

A <sub>3</sub>	A <sub>2</sub>	A <sub>1</sub>	A <sub>0</sub>	y
0	0	0	0	0
0	0	0	1	0
0	0	1	0	0
0	0	1	1	1
0	1	0	0	1
0	1	0	1	1
0	1	1	0	1
0	1	1	1	1
1	0	0	0	0
1	0	0	1	0
1	0	1	0	0
1	0	1	1	0
1	1	0	0	0
1	1	0	1	0
1	1	1	0	0
1	1	1	1	0

k-map simplification

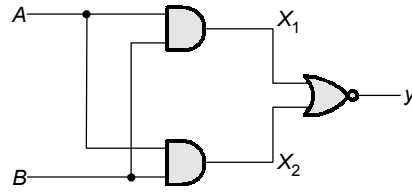


$$y = \bar{A}_3A_2 + \bar{A}_3A_1A_0 = \bar{A}_3(A_2 + A_1A_0)$$

12. (a)

$$\begin{aligned} x &= (M + N)(\bar{M} + P)(\bar{N} + \bar{P}) \\ &= (M\bar{M} + MP + N\bar{M} + NP)(\bar{N} + \bar{P}) \\ &= (MP + N\bar{M} + NP)(\bar{N} + \bar{P}) \\ &= MP\bar{N} + MP\bar{P} + N\bar{M}\bar{N} + N\bar{M}\bar{P} + NP\bar{N} + NP\bar{P} \\ x &= MP\bar{N} + N\bar{M}\bar{P} \end{aligned}$$

13. (d)



$$X_1 = AB$$

$$X_2 = AB$$

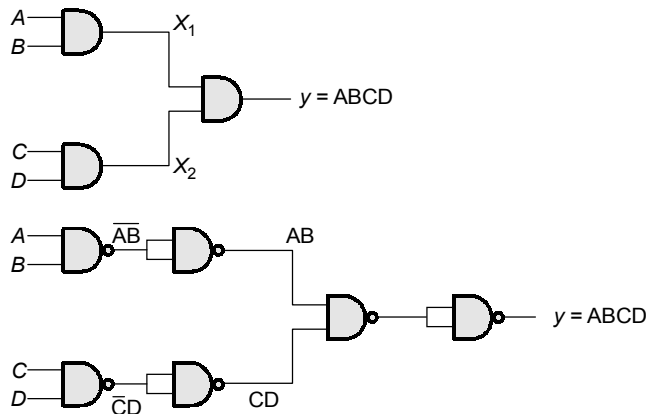
$$y = \overline{AB + AB} = \overline{AB}$$

A	B	$y = \overline{AB}$
0	0	1
0	1	1
1	0	1
1	1	0

→ LED will be ON

14. (c)

Implementing  $y$  using 2 input AND gates



∴ Hence total 6, 2 input NAND gates are required.

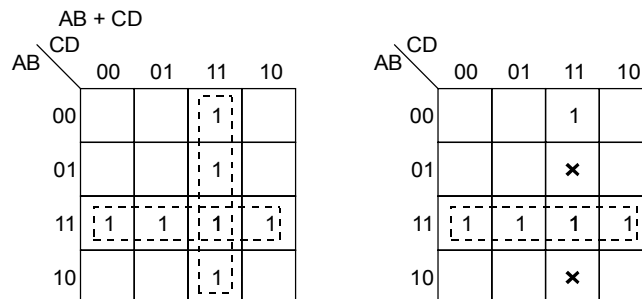
15. (c)

The expression of output is

$$r = (\bar{A} + B)(A + \bar{B}) = A\bar{A} + \bar{A}\bar{B} + AB + B\bar{B} = AB + \bar{A}\bar{B} = A \odot B$$

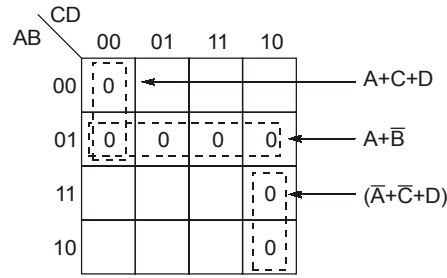
16. (a)

k-map



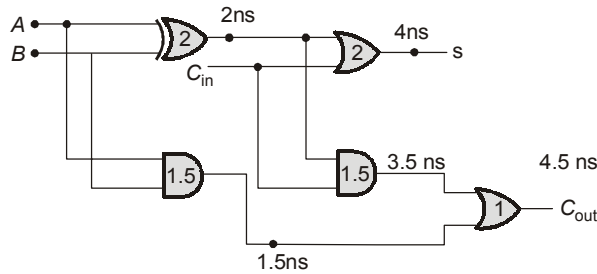
Comparing both the k-maps, it is clear that there are two don't care entries. Which are (7, 11).

17. (a)  
Drawing k-map for x.



∴ 
$$x = (A + C + D)(\bar{A} + \bar{C} + D)(A + \bar{B})$$

18. (a)  
Full adder implementation

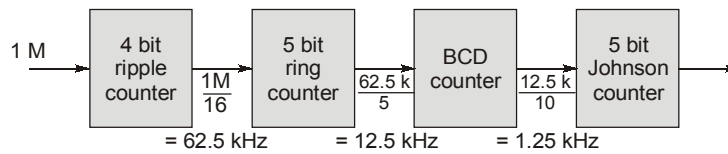


$$s = A \oplus B \oplus C$$

$$C_{out} = (A \oplus B) C_{in} + AB$$

From figure, it is clear that delay for sum and carry will be 4ns and 4.5ns respectively.

19. (c)  
4 bit ripple ripple counter has  $2^4 = 16$  states  
5 bit ring counter has 5 states  
BCD counter has 10 states  
5 bit Johnson counter has  $5 \times 2 = 10$  states



$$FOS = \frac{1250}{10} = 125 \text{ Hz}$$

21. (d)  
Propagation delay

$$T_d = (n - 1) t_{p \text{ carry}} + \max(t_{p \text{ carry}}, t_{p \text{ sum}})$$

$$= (16 - 1) \times 5 + \max(5, 6)$$

$$= 75 + 6 = 81 \text{ ns}$$

## 22. (a)

For given quadratic equation,

$$(x - 3) = 0$$

and

$$(x - 6) = 0 \text{ will be two factors}$$

$\therefore$

$$(x - 3)(x - 6) = x^2 - 11x + 22 \text{ both are equal}$$

$\therefore$

$$x^2 - (6 + 3)x + 6 \times 3 = x^2 - 11x + 22$$

Let  $b$  be the base of number

$$6 + 3 = 1b + 1$$

$$b = 8$$

Also,

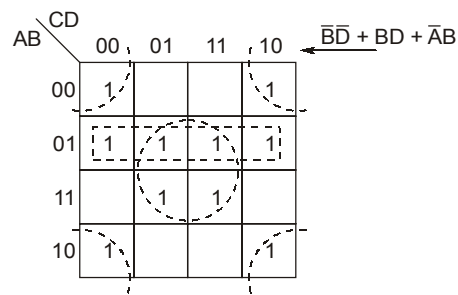
$$6 \times 3 = 2b + 2$$

$$18 = 2b + 2$$

$$b = 8$$

## 23. (a)

$k$  map



## 24. (b)

Let  $n$  number of flip flop cascaded each having propagation delay of  $t_{pd}$ .

Frequency of operation

$$\frac{1}{nt_{pd}} \geq 10 \text{ MHz}$$

$$\therefore n \leq \frac{1}{t_{pd} \times 10 \text{ MHz}} \leq \frac{1}{12 \times 10^{-9} \times 10^7} \leq \frac{100}{12}$$

$$n = 8$$

For  $n = 8$  MOD number of counter  $2^8 = 256$

## 25. (a)

When odd number of NOT gate are arranged with output as a feedback to input it acts as an oscillator with frequency,

$$f = \frac{1}{2nt_{pd}} \quad n = \text{Number of NOT gates}$$

$$\therefore f = \frac{1}{2 \times 5 \times t_{pd}}$$

$$\therefore t_{pd} = \frac{1}{10 \times f_{\max}} = \frac{1}{10 \times 20} \mu\text{s} = \frac{1000}{10 \times 20} \mu\text{s} = 5 \text{ ns}$$

$\therefore$  Each NOT gate has a propagation delay of 5 ns.

26. (a)

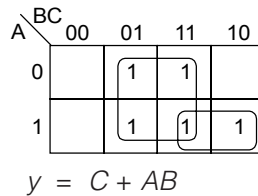
Expression of  $E$ .

$$E = \bar{A}O + A \cdot B = AB$$

$$y = \bar{E}C + E \cdot 1 = \overline{ABC} + AB$$

$$= (\bar{A} + \bar{B})C + AB = \bar{A}C + \bar{B}C + AB$$

k-map simplification



27. (a)

An eight bit binary ripple counter will be RESET i.e. initial condition 00000000 after

$$= (\text{Mod number} - \text{Current state})$$

$$= 256 - 127 \quad (0111 \ 1111 = 127_{10})$$

$$= 129 \text{ clock pulses}$$

$\therefore$  i.e. after 129 clock pulses counter will be at  $(00)_{10}$ .

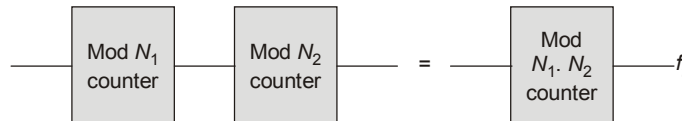
After additional  $(135 - 129) = 6$  pulses counter will be at 0000 0110 state.

28. (c)

(i)  $n$  bit Johnson counter will have  $2n$  states

$\therefore$  3bit Johnson counter has mod number =  $3 \times 2 = 6$

In 2nd stage, D-Flop acts as a T-flip flop and has mod no. 2



Hence given counter arrangement has  $6 \times 2 = 12$  mod number

29. (a)

$$\text{Full scale voltage} = \text{Step size} \times \text{Maximum count}$$

$$= \text{Step size} \times (2^n - 1)$$

$$= 10 \text{ mV} \times 255$$

$$= 2.55 \text{ V}$$

$$\% \text{ Resolution} = \frac{1}{2^n - 1} \times 100 = \frac{100}{255} = 0.392\%$$

30. (d)

Number of function using  $n$  variables is  $2^{2^n}$ .

Since, here  $n = 3$ , so number of function is 256.

