

# CLASS TEST

S.No. : 07 BS\_CS\_B\_110919

Digital Logic



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# CLASS TEST 2019-2020

## COMPUTER SCIENCE & IT

Date of Test : 11/09/2019

### ANSWER KEY > Digital Logic

1. (b)	7. (c)	13. (d)	19. (d)	25. (d)
2. (a)	8. (c)	14. (b)	20. (c)	26. (c)
3. (b)	9. (c)	15. (b)	21. (d)	27. (b)
4. (d)	10. (a)	16. (a)	22. (c)	28. (b)
5. (b)	11. (a)	17. (d)	23. (c)	29. (b)
6. (c)	12. (c)	18. (c)	24. (b)	30. (b)

**DETAILED EXPLANATIONS**

1. (b)

For every combination of x, y, z the function value remains same for input  $\bar{x}, y, \bar{z}$ .

x	y	z	$f(x, y, z) = f(\bar{x}, y, \bar{z})$
0	0	0	} either 0 or 1
1	0	1	
0	0	1	} either 0 or 1
1	0	0	
0	1	0	} either 0 or 1
1	1	1	
0	1	1	} either 0 or 1
1	1	0	

Effectively there are only four rows for the truth table of the function  $f(x, y, z)$ .

$\therefore$  Total Boolean expressions possible is  $2^4 = 16$ .

2. (a)

$$(135)_x + (144)_x = (323)_x$$

$$\Rightarrow (x^2 + 3x + 5) + (x^2 + 4x + 4) = 3x^2 + 2x + 3$$

$$\Rightarrow x^2 - 5x - 6 = 0$$

$$\Rightarrow (x - 6)(x + 1) = 0$$

$$\Rightarrow x = 6$$

[ $\therefore$  base of a number cannot be -ve]

3. (b)

Let initial  $Q_1 = Q_0 = 0$

$J_0$	$K_0$	$J_1$	$K_1$	$Q_1$	$Q_0$
				0	0
1	1	0	1	0	1
1	1	1	1	1	0
0	1	0	1	0	0

It is a mod-3 counter ( $00 \rightarrow 01 \rightarrow 10 \rightarrow 00$ )

So,  $K = 3$

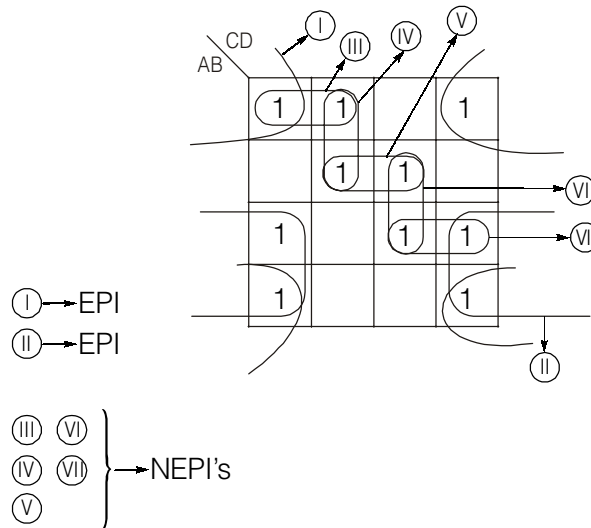
4. (d)

$$37 = 16 \times 2 + 5$$

After 37 clock pulses, the state of MOD-16 DOWN counter will be five states below the present state.

$$\begin{array}{r} 0110 (6) \\ -0101 (-5) \\ \hline 0001 \end{array}$$

5. (b)

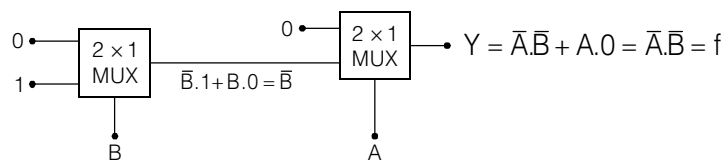


EPI = Essential Prime Implicant [which cover a minterm not covered by any other term]  
NEPI = Non Essential Prime Implicant. Number of EPI's = 2, number of NEPI's = 5.

6. (c)

$$f = \bar{A}\bar{B}C + \bar{A}\bar{B}\bar{C} = \bar{A}\bar{B}(C + \bar{C}) = \bar{A}\bar{B}$$

with  $2 \times 1$  multiplexer



7. (c)

Range of signed 1's complement number is  $-2^{n-1} + 1$  to  $2^{n-1} - 1$ .

8. (c)

$$\begin{aligned} (AB + B\bar{C} + \bar{A}C)(A + C) &= AB + AB\bar{C} + ABC + \bar{A}C \\ &= AB(1 + \bar{C} + C) + \bar{A}C \\ &= AB + \bar{A}C \end{aligned}$$

9. (c)

Multiplexer is a combinational circuit, converts parallel to serial data and it selects one of the several inputs and transmits to a single output.

10. (a)

$$\begin{aligned} (x \oplus y) &= \bar{x}y + x\bar{y} \\ &= \overline{xy + \bar{x}\bar{y}} \\ &= \overline{0 + \bar{x}\bar{y}} \quad (\text{as } xy = 0) \\ &= (x + y) \end{aligned}$$

11. (a)

$$(36)_7 = 7 \times 3 + 6 = (27)_{10}$$

$$(67)_8 = 8 \times 6 + 7 = (55)_{10}$$

$$(98)_{10} = 9 \times 10 + 8 = (98)_{10}$$

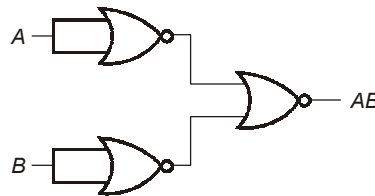
$$(34)_5 = 3 \times 5 + 4 = (19)_{10}$$

$$\therefore 27 + 55 + 98 + 19 = (199)_{10}$$

$$\begin{array}{r} 9 \overline{)199} 1 \\ 9 \overline{)22} 4 \\ \underline{2} \end{array} \Rightarrow (199)_{10} = (241)_9 = 241$$

12. (c)

$$\begin{aligned} f &= \overline{\overline{\overline{A + [B + \overline{C(AB + AC)]}}}]} \\ &= A \cdot [B + \overline{C(AB + AC)}] && \text{[Demorgan's Law]} \\ &= A \cdot [B + \overline{C}(\overline{AB} \cdot \overline{AC})] && \text{[Demorgan's Law]} \\ &= A \cdot [B + \overline{C}(\overline{A} + \overline{B})(\overline{A} + \overline{C})] && \text{[Demorgan's Law]} \\ &= A[B + \overline{C}(\overline{A} + \overline{BC})] && \text{[Distributive property]} \\ &= A[B + \overline{A}\overline{C}] = AB + A\overline{A}\overline{C} \\ &= AB \text{ (AND gate to be implemented)} \end{aligned}$$



⇒ Minimum number of NOR gate required = 3

13. (d)

Clock	D	$Q_D$	T	$Q_T$
-	0	0	1	1
1	1	0	1	0
2	1	1	0	1
3	1	1	0	1
4	1	1	0	1

The circuit will be locked up in the state '11'. So state of circuit after 7 clock pulses is '11'.

14. (b)

$$T \geq 3 \times 50 \text{ n sec}$$

$$T \geq 150 \text{ n sec}$$

$$T_{\min} = 150 \text{ nsec}$$

$$f_{\max} = \frac{1 \times 10^9}{150} = 6.67 \text{ MHz}$$

15. (b)

Total AND gates for a n-bit carry look ahead generator is  $1 + 2 + 3 + \dots + n = \frac{n(n+1)}{2}$

Here  $n = 5 \Rightarrow \frac{5(5+1)}{2} = 5 \times 3 = 15$

16. (a)

Let us consider active high input

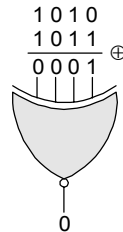
		YZ			
X		00	01	11	10
0		0	1	1	0
1		0	1	0	1

$$F = \Sigma(1, 3, 5, 6) = \Pi(0, 2, 4, 7)$$

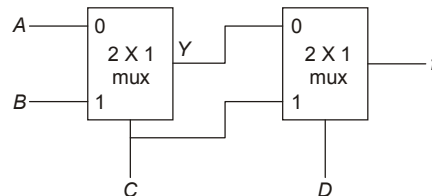
$$= (Y + Z) \cdot (X + Z) \cdot (\bar{X} + \bar{Y} + \bar{Z})$$

17. (d)

In EX-NOR gate, if odd number of inputs are 1 then output is zero.



18. (c)



$$Y = A\bar{C} + BC$$

$$f = Y\bar{D} + CD$$

$$f = (A\bar{C} + BC)\bar{D} + CD = A\bar{C}\bar{D} + BC\bar{D} + CD$$

$$\begin{aligned}
 &= A(B + \bar{B})\bar{C}\bar{D} + (A + \bar{A})CBC\bar{D} + CD(A + \bar{A})(B + \bar{B}) \\
 &= AB\bar{C}\bar{D} + A\bar{B}\bar{C}\bar{D} + ABC\bar{D} + \bar{A}BC\bar{D} + ABCD + A\bar{B}CD + \bar{A}BCD + \bar{A}\bar{B}CD \\
 f(A, B, C, D) &= \Sigma m(3, 6, 7, 8, 11, 12, 14, 15)
 \end{aligned}$$

		CD	
AB			1
			1
	1		1
	1		1

19. (d)

$$\begin{aligned}
 f_1(A, B, C) &= \Sigma(2, 3, 4) \\
 f_2(A, B, C) &= \pi(0, 1, 3, 6, 7) = \Sigma(2, 4, 5)
 \end{aligned}$$

For function  $f$  to be zero:

$$\begin{aligned}
 f_3(A, B, C) &= \overline{[f_1(A, B, C) \cap f_2(A, B, C)]} \\
 &= \Sigma(0, 1, 3, 5, 6, 7)
 \end{aligned}$$

Maximum minterms possible are 6.

20. (c)

$$\text{Number of level-1 MUX} = \frac{4}{2} = 2$$

$$\Rightarrow \text{Number of level-2 MUX} = \frac{2}{2} = 1$$

$$\therefore \text{Total MUX required} = 3$$

21. (d)

$$\begin{array}{r}
 A = \quad \quad \quad a_2 \quad \quad a_1 \quad \quad a_0 \\
 B = \quad \quad \quad \quad \quad \quad b_1 \quad \quad b_0 \\
 \hline
 A \times B = \quad \quad a_2b_0 \quad a_1b_0 \quad a_0b_0 \\
 \quad \quad \quad b_1a_2 \quad b_1a_1 \quad b_1a_0 \quad \downarrow \\
 \hline
 \quad \quad \quad b_1a_2 \quad (a_2b_0 + a_1b_1) \quad (a_1b_0 + b_1a_0) \quad a_0b_0 \\
 \quad \quad \quad C_3 \quad \quad C_2 \quad \quad C_1 \quad \quad C_0
 \end{array}$$

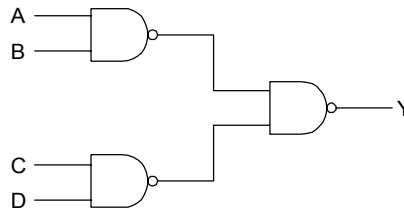
Number of AND gates required  $X = 6$

Number of one bit full adders required  $Y = 3$

$$X + Y = 6 + 3 = 9$$

22. (c)

$$Y = AB + CD = \overline{\overline{AB}} + \overline{\overline{CD}}$$



23. (c)

Let the base of the system is B

$$9 \times 8 \times 5 = (190)_B$$

$$(360)_{10} = (190)_B$$

$$360 = B^2 + 9B + 0$$

$$B^2 + 9B - 360 = 0$$

$$B = 15,$$

$$B = -24$$

Since base can not be negative.

⇒ Base

$$B = 15$$

24. (b)

$$\text{Decimal input} = 92$$

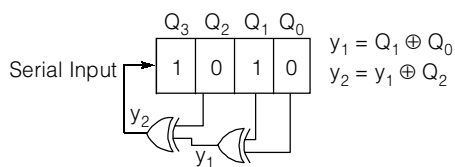
$$\text{BCD} = 10010010$$

$$\text{Output of Gray code converter} = 11011011$$

$$Y_0 \text{ corresponds to } I_m \text{ with } (S_n, \dots, S_0) \text{ is } = (11011011)_2$$

$$m = 219$$

25. (d)



CLK pulse		Q <sub>3</sub>	Q <sub>2</sub>	Q <sub>1</sub>	Q <sub>0</sub>
1	0	1	0	1	0
2	1	1	1	0	1
3	2	0	1	1	0
4	3	0	0	1	1
5	4	0	0	0	1
6	5	1	0	0	0
7	6	0	1	0	0
7	7	1	0	1	0

$$Q_3(t + 1) = Q_0(t) \oplus Q_1(t) \oplus Q_2(t)$$

26. (c)

Consider characteristic equation of J-K Flip-Flop:  $Q_{2N+1} = J\bar{Q}_{2N} + \bar{K}Q_{2N}$

$$J = \bar{Q}_0; K = 1 \Rightarrow Q_{2N+1} = \bar{Q}_0 \bar{Q}_{2N} \quad \dots(i)$$

$$\text{if } Q_0 = 1 \Rightarrow Q_{2N+1} = 0$$

$$\text{if } Q_0 = 0 \Rightarrow Q_{2N+1} = \bar{Q}_{2N}$$

Consider characteristic equation of D – Flip-Flop:  $Q_{N+1} = D$

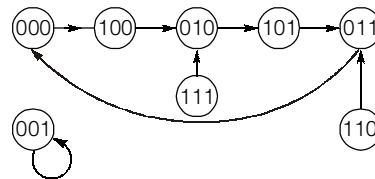
$$Q_{N+1} = Q_2 \quad \dots(ii)$$

Consider characteristic equation of T – Flip-Flop:  $Q_{N+1} = T \oplus Q_N$

$$Q_{N+1} = Q_1 \oplus Q_0 \quad \dots(iii)$$

Using equations (i), (ii) and (iii)

Present State			Next State		
$Q_2$	$Q_1$	$Q_0$	$Q_2$	$Q_1$	$Q_0$
0	0	0	1	0	0
0	0	1	0	0	1
0	1	0	1	0	1
0	1	1	0	0	0
1	0	0	0	1	0
1	0	1	0	1	1
1	1	0	0	1	1
1	1	1	0	1	0



The number of used states = 5

∴ Modulus value of the counter = 5

27. (b)

$$A = (EC)_{16} \rightarrow (1032)_6$$

$$B = (20)_{10} \rightarrow (32)_6$$

$$C = (15)_8 \rightarrow (21)_6$$

$$S = A + B + C$$

$$= (1085)_{10}$$

(add in base-6 and convert into base-10)



28. (b)

To obtain a JK flip-flop from a T flip-flop, we first construct the characteristic table of JK flip-flop; and then obtain the excitation values for the T flip-flop as shown below:

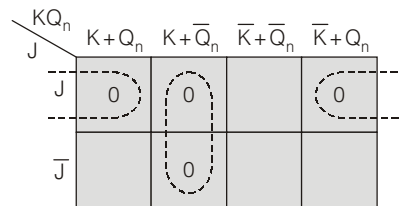
J	K	$Q_n$	$Q_{n+1}$	T
0	0	0	0	0
0	0	1	1	0
0	1	0	0	0
0	1	1	0	1
1	0	0	1	1
1	0	1	1	0
1	1	0	1	1
1	1	1	0	1

Characteristics table of JK flip flop ←

→ Excitation values for T flip-flop

Now, assuming T to be an output, we solve it in terms of J, K,  $Q_n$  inputs. This gives the definition of the logic to be applied on the T input.

Also, observing the given options, we solve for T using a maxterms map instead of using a minterms map, as shown below:



$$T = (J + Q_n) \cdot (K + \bar{Q}_n)$$

The circuit corresponding to this expression is given option (b).

29. (b)

$$F = \bar{A}\bar{B}C + \bar{A}BC + A\bar{B}\bar{C} + ABC$$

$$F = \bar{A}C(\bar{B} + B) + A\bar{C}(\bar{B} + B)$$

$$F = \bar{A}C + A\bar{C}$$

$$F = A \oplus C$$

30. (b)

$$D_1 = Q_0, D_0 = Q_0 \oplus Q_1$$

$$\Rightarrow (Q_1 Q_0) 00 \rightarrow 01 \rightarrow 10 \rightarrow 00 \rightarrow 01 \dots$$

